

# A Logic of Isolation

Can Başkent, David Gilbert and Giorgio Venturi

Başkent: Department of Computer Science, Middlesex University, England Gilbert: Department of Philosophy, University of British Columbia, Canada Venturi: Dipartimento di Civilitá e Forme del Sapere, University of Pisa, Italy

c.baskent@mdx.ac.uk canbaskent.net/logic

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## Topological isolation operator

## enjoys the richness of topologies

## with a well-defined modality,

## and can be used to express "factive ignorance".

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# Introduction

Topological semantics remains the oldest semantics for modal logics. Modal operators have traditionally been identified with interior and closure operators, as well as, less often so, with border and boundary operators. The tuple  $(S, \sigma)$  is topological space for set S and the collection  $\sigma$  of subsets of S satisfying the following axioms:

- 1. The empty set and S are in  $\sigma$
- 2.  $\sigma$  is closed under arbitrary unions,
- 3.  $\sigma$  is closed under finite intersections.

The collection  $\sigma$  is called a **topology**. We call the elements of  $\sigma$  **opens**, and their complements **closeds**.

The largest open set contained in a set is called the **interior**. The smallest closed set containing a given set is called the **closure**.

The interior operator Int returns the interior of a given set. Dually, the closure operator Clo returns the closure of a given set.

In modal logic, Int is identified with the classical modal operator  $\Box$ . Dually and similarly, Clo is identified with  $\Diamond$ .

In a topological model  $M = (S, \sigma, V)$ , where V is a valuation function in the ordinary sense, the semantics of the topological modal operators is given as follows, for  $w \in S$ .

$$\begin{array}{ll} M, w \models \Box \varphi & iff \quad \exists U \in \sigma \ (w \in U \land \forall v \in U . M, v \models \varphi) \\ M, w \models \Diamond \varphi & iff \quad \forall U \in \sigma \ (w \in U \rightarrow \exists v \in U . M, v \models \varphi) \end{array}$$

Notice the difference in complexity of the topological semantics for modalities!

## Basics of Topological Modal Models iv

Under these assumptions, McKinsey and Tarski showed in 1944 that the modal logic of topological spaces is S4 by identifying the axioms of S4 modal logic with the axioms of the closure/interior operator.

S4 is the logic of reflexive and transitive frames, characterised by the following axioms.

- All axioms of propositional logic
- $\cdot \ \Box(\varphi \to \psi) \to (\Box \varphi \to \Box \psi)$
- $\boldsymbol{\cdot} \ \Box \varphi \to \varphi$
- $\cdot \ \Box \varphi \to \Box \Box \varphi$

#### Reference

J. C. C. McKinsey and A. Tarski, "The algebra of topology", *Annals of Mathematics*, vol. 45(1), pp. 141–191, 1944.

**Isolated Points** 

Given a topological space  $(S, \sigma)$  and  $X \subseteq S$ , the point x is an **isolated** point of X if there is an open neighbourhood U of x such that  $U \cap S = \{x\}$ .

Here is an illustration from Wikipedia, showing that 0 is an isolated point in the set  $[1, 2] \cup \{0\}$ , with natural Euclidean topology:



# A point x is isolated in space $(S, \sigma)$ if and only if $\{x\}$ is open. In $\mathbb{N}$ , each natural number $n \in \mathbb{N}$ is isolated.

# Logic of Isolated Points

Factive ignorance postulates that

If an agent is ignorant of  $\varphi$ , then  $\varphi$  is true.

In a classical Kripke semantics, its semantics is given as follows for a model  $M = \langle W, R, V \rangle$ 

 $\begin{array}{ll} M,w\models \mathsf{I}\varphi & \textit{if and only if} & \forall w'\neq w \textit{ and such that }\mathsf{Rww'},\\ M,w'\not\models\varphi \textit{ and }M,w\models\varphi \end{array}$ 

where I $\varphi$  represents that "the agent is ignorant that  $\varphi$ ".

#### Reference

Kubyshkina, E. and Petrolo, M, "A logic for factive ignorance", *Synthese*, vol. 198, pp. 5917–5928, 2021.

#### The syntax of Logic of Isolation is as follows:

$$\alpha ::= p \mid \neg \alpha \mid \alpha \land \alpha \mid [\mathbf{i}]\alpha$$

The axioms and the proof rules of Logic of Isolation are given as follows:

- All instances of propositional tautologies,
- ·  $[i] \varphi o \varphi$ ,
- · ([i] $\varphi \wedge$  [i] $\psi$ )  $\rightarrow$  [i]( $\varphi \wedge \psi$ ),
- From  $\vdash \varphi$  and  $\vdash \varphi \rightarrow \psi$  infer  $\vdash \psi$ ,
- From  $\vdash \varphi \rightarrow \psi$  infer  $\vdash \varphi \rightarrow ([i]\psi \rightarrow [i]\varphi)$

We call this system  $S^{i}$ .

The following are all theorems of  $S^i$ .

- $([i]\varphi \wedge [i]\psi) \rightarrow [i](\varphi \wedge \psi)$
- $[i](\varphi \lor \psi) \to ([i]\varphi \lor [i]\psi)$
- ·  $[i]\varphi \rightarrow [i][i]\varphi$

One can also derive  $\vdash$  [i] $\varphi \leftrightarrow$  [i] $\psi$  from  $\vdash \varphi \leftrightarrow \psi$ .

## Logic of Isolation: Semantics

Relational modal semantics for  $S^{i}$  is given as follows:

 $M, w \models [i]\varphi$  iff  $(M, w \models \varphi \text{ and } \forall w' \neq w (wRw' \text{ implies } M, w' \not\models \varphi))$ 

#### Theorem

**S**<sup>i</sup> is sound and strongly complete with respect to the class of all relational frames.

#### Reference

D. Gilbert, E. Kubyshkina, M. Petrolo, and G. Venturi, "Logics of Ignorance and Being Wrong", *Logic Journal of the IGPL*, vol.30(5), pp. 870–885, 2021.

# Topological Semantics for Logic of Isolation

A neighborhood frame is a pair (X, N) such that  $X \neq \emptyset$  and  $N : X \mapsto \wp(\wp(X))$ .

A neighborhood model is a pair  $\langle F, V \rangle$ , where F is a neighborhood frame and V : Prop  $\mapsto \wp(X)$  is a valuation function.

Standard topological definition of isolation can be extended to neighborhood systems, where less mathematical structure is insisted upon:

x is an isolated point of S if there is a neighborhood U of x, i.e.,  $U \in N(x)$ , such that  $U \cap S = \{x\}$ .

Given a model *M* and a formula  $\alpha$ , the *truth set* of  $\alpha$  in *M* is denoted by  $[\![\alpha]\!]^M$ . We will omit the superscript for easy reading.

The language of  $\mathbf{S}^{i}$  is reflexive-insensitive.

The satisfaction of formulas in the language of  $S^i$  in a model  $M = \langle W, R, V \rangle$  is not affected when arbitrary elements of the form (w, w), where  $w \in W$ , are either added to or removed from R.

Particularly, because the definition of  $[\cdot]$  utilizes only sets of each N(x) that contain x, the addition or removal of sets that do not contain x will be immaterial.

For a given neighborhood frame  $\langle X, N \rangle$ , consider the set

 $S_{x} := \{Y \in \wp(X) : x \notin Y\}$ 

Given a neighborhood frame  $F = \langle X, N \rangle$ , construct the frames  $F = \langle X, N^+ \rangle$  and  $F = \langle X, N^- \rangle$ , where  $N^+$  and  $N^-$  are defined as follows for all  $x \in X$ :

 $N^+(x) := N(x) \cup S_x$ 

 $N^{-}(x) := N(x) \setminus S_x$ 

If 
$$M = \langle X, N, V \rangle$$
, then  $M^+ = \langle X, N^+, V \rangle$  and  $M^- = \langle X, N^-, V \rangle$ .

#### Theorem

Let M be a neighborhood model. Then M, M<sup>+</sup>, and M<sup>-</sup> (as well as the intermediate models) are all pointwise equivalent. That is,

$$\llbracket \alpha \rrbracket^{\mathsf{M}^-} = \llbracket \alpha \rrbracket^{\mathsf{M}} = \llbracket \alpha \rrbracket^{\mathsf{M}^+}$$

for all formulas  $\alpha$  in the language of **S**<sup>*i*</sup>.

**Supplemented Neighbourhood Models** A neighborhood frame is *supplemented* when its neighborhood function is closed under supersets: for every x, if  $Y \subseteq N(x)$  and  $Y \subseteq Z$ , then  $Z \subseteq N(x)$ .

Given a neighborhood frame,  $F = \langle X, N \rangle$ , let  $F^s = \langle X, N^s \rangle$  be the supplementation of F when, for all  $x \in X$ :

$$N^{s}(x) = \{Y \subseteq \wp(X) : \exists U \subseteq Y \text{ s.t. } U \in N(x)\}$$

For a model  $M = \langle F, V \rangle$ , let  $M^{s} = \langle F^{s}, V \rangle$ .

Anchored Neighborhood System A neighborhood function (and, hence, the resulting system) is *anchored* when, for every point  $x \in X$ ,

 $\forall U \in N(x) (x \in U)$ 

Note that we do not force  $N(x) \neq \emptyset$  in order to be anchored.

Theorem

Let M be an anchored neighborhood model. Then

 $\llbracket \alpha \rrbracket^{\mathsf{M}} = \llbracket \alpha \rrbracket^{\mathsf{M}^{\mathsf{s}}}$ 

#### Theorem

**S**<sup>i</sup> is sound with respect to the class of neighborhood frames that are closed under intersections.

#### Theorem

**S**<sup>i</sup> is strongly complete with respect to the class of neighborhood frames that are anchored and closed under intersections.

#### Theorem

**S**<sup>i</sup> is strongly complete with respect to the class of neighborhood frames that are anchored, closed under intersections, and supplemented.

#### Theorem

**S**<sup>i</sup> is sound and (strongly) complete with respect to the class of all augmented neighborhood frames and all anchored, augmented neighborhood frames.

A neighborhood system is *discrete* when  $\{x\} \in N(x)$ , for every  $x \in X$ . Consider the following axiom schema:

$$\varphi \leftrightarrow [i]\varphi$$
 (Disc)

Call **SDisc** the system obtained by adding (Disc) to  $S^{i}$ .

In the presence of (Disc), no other modal axioms are necessary and neither is the rule "From  $\vdash \varphi \rightarrow \psi$  infer  $\vdash \varphi \rightarrow ([i]\psi \rightarrow [i]\varphi)$ ".

SDisc can be axiomatized by the following:

- All instances of propositional tautologies,
- ·  $\varphi \leftrightarrow [\mathbf{i}] \varphi$ ,
- From  $\vdash \varphi$  and  $\vdash \varphi \rightarrow \psi$  infer  $\vdash \psi$

#### Theorem

**SDisc** is valid on a relational frame F if and only if, for each w, if wRz, then w = z.

#### Theorem

**SDisc** is sound and strongly complete with respect to the class of relational frames in which, for each w, if wRz, then w = z.

#### Theorem

**SDisc** is sound and strongly complete with respect to the class of discrete neighborhood systems.

(Hence, due to the semantic insensitivities noted above, also with respect to anchored, discrete, supplemented neighborhood systems.)

Let us now transition to topologies proper.

Satisfaction at points in a topo-model is defined as usual, with the clause for [i] resembling closely the one given for neighborhood systems, but with reference to the topology  $\sigma$  rather than the neighborhood function *N*:

$$\llbracket [i]\varphi \rrbracket := \{ x : \exists U \in \sigma \text{ s.t. } U \cap \llbracket \varphi \rrbracket = \{ x \} \}$$

With the semantics so defined,  $\mathbf{S}^{i}$  is sound with respect to the class of all topo-models.

Moreover, **SDisc** is sound with respect to the class of all discrete topological spaces, since all singletons are open (and closed).

Theorem

**SDisc** is complete with respect to the class of all discrete topological spaces.

Conclusion

## Conclusion

We presented a geometric interpretation of [i] as an isolated points operator in a variety of systems.

The next question is whether there exist intermediate logics (between **S**<sup>i</sup> and **SDisc**) that characterize interesting classes of either neighborhood systems or topologies.

The extensions of  $S^i$  with only neighborhood systems in mind is an interesting direction to pursue. For example, one can consider adding to  $S^i$  the axiom  $\neg[i]\top$ .

Relationally, this forces all worlds to be non-reflexively serial (for each x, there is a  $y \neq x$  such that xRy). This axiom forces a lack of discreteness (hence, the resulting logic is not an intermediate logic, but inconsistent with **SDisc**), and that it is sound and strongly complete with respect to the class of all neighborhood systems in which  $\{x\} \notin N(x)$ .

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## Topological isolation operator

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# Thank you!

Talk slides are available at my website

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